# $\int_{0}^{1} f=\mathbb{E}(f(\mathcal{C}(\mathcal{G}(\varepsilon))))$ <br> Integration is Nothing but a Discrete Sum 

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#### Abstract

In this short article we will introduce Riemann Integration of a function on $(0,1)$ using the Collatz Bijection $\mathcal{C}$, which claims that there exists some sort of bridge between natural integers and real numbers in $(0,1)$. The purpose of the Geometric Law $\mathcal{G}(\varepsilon)$ with a small probability of success $\varepsilon$ is to mimic the limit uniform law on $\mathbb{N}$. The main result of the article is as follows.


Let $f$ be Riemann Integrable on $(0,1)$, let $D>0$ be the precision parameter:

$$
\exists \varepsilon>0, \quad\left|\int_{0}^{1} f(x) \mathrm{d} x-\sum_{n=1}^{\infty} \varepsilon f(\mathcal{C}(n))(1-\varepsilon)^{n}\right|<2^{-D}
$$

With an upper bound on $\varepsilon$ under certain conditions.
On a more philosophical note, it means that continuous analysis is discrete, and even finite if you neglect the remainder, and only consider an approximation. A higher result would be to claim, that no matter $D$, $\exists \varepsilon, N>0$ such that:

$$
\left|\int_{0}^{1} f(x) \mathrm{d} x-\sum_{n=1}^{N} \varepsilon f(\mathcal{C}(n))(1-\varepsilon)^{n}\right|<2^{-D}
$$

## 1 Introduction to my understanding of Riemann Integration

The idea behind Riemann integration lies in a very elegant formula.

$$
\inf _{n \geq 1}\left\{\left|\int_{0}^{1} f(x) \mathrm{d} x-\frac{1}{n} \sum_{k=0}^{n-1} f(k / n)\right|\right\}=0
$$

Which implies that an integral from 0 to 1 of a function, is none but the average of its values. The problem with working on the average is that, you sometimes need a lot of samples to generate a coherent value, yet, again, the problem of having an asymptotical value at $k / n$. Making the series hardly integrated on say $\int_{0}^{1} P(x)^{-2} \mathrm{~d} x$ with lots of roots in $\mathbb{Q} \cap(0,1)$. There are actual tricks to challenge this method of integration. But it needs knowledge on the asymptotical behavior of the function, in order to consider it a Dirac, which ultimately leads to the Riemann Integration as the average of non infinite terms, plus a Dirac coefficients, collection.

## 2 Introduction to my Understanding of Measure Theory

I know probabilities are measure, and I also know that the measure $\mu_{0}(x)=$ $(x \in \mathbb{N})$ exists, but cannot be interpreted as a probability measurement for the simple reason the $\mathbb{N}$ is the discrete infinite $\aleph$. Basically, suppose you have finite ressources, and want every integers to have an equal part, it has to be 0 , which makes no sense because the series of 0 is zero, and not 1 as would discrete probability suggest. So let's make a thought experiment and try to imagine what should be the closest thing to this uniform measurement might be. You give people a lot, in fact everything on the probability that the person succeeded an impossible task, and that no previous person already made it. We call that probability $\varepsilon$, you can think of it as winning the Million dollar, every week for 52 consecutive weeks. Two things to know about it, it's possible, and it is very small. $\mu(x)=(x \in \mathbb{N}) \varepsilon(1-\varepsilon)^{x}$, this is what we call a Geometric Law, of probability of success $\varepsilon$ and the closest thing we know to communism on the premise that lottery ticket were free. In fact, you can prove that:

$$
\mu_{0}-\frac{\mu}{\varepsilon}
$$

Is in physics what we call a $o(1)$.

## 3 Introduction to my understanding of the Collatz Conjecture

It all starts with a very simple but elegant equation.

$$
C(x)=\frac{1}{2}((x \equiv 1[2])(3 x+1)+(x \equiv 0[2]) x)
$$

And then a twisted mind of a scientist who spent way too long documenting the subject to claim that:

$$
\mathcal{C}(x)=\frac{1}{2} \sum_{k=0}^{\infty}\left(C^{k}(x) \equiv 1[2]\right) 2^{-k}
$$

Is both:

- Defined on $\mathbb{N}$, natural integers, zero being zero.
- A Bijection from $\mathbb{N}$ to $[0,1)$ which can be easily disproven if you can prove Collatz conjecture to be True.

It is what we call a computational bijection with very interesting properties.

## 4 How it is all related

It can be summarized in simple terms: $\int_{0}^{1} f(x) \mathrm{d} x$ is a measurement $\mu_{1}(f)$ evaluated in a continuous motion. But that continuous motion, can be actually emulated by taking random samples uniform and increase the number of samples. It suffice to prove that $\mathcal{C}(\mathcal{G}(\varepsilon))$ converges to $\mathcal{U}(0,1)$. The core of the proof is that for fixed arbitrary $a<b$.

$$
\operatorname{Pr}(a \leq X \leq b)=b-a+O(\varepsilon)
$$

Which is trivial to prove, when you know that:

- $\forall a<b, C^{a}\left(x+2^{b}\right) \equiv C^{a}(x)[2]$
- $C^{a}\left(x+2^{a}\right) \equiv 1+C^{a}(x)[2]$

There is a bijection between $\left\{1,2,3, \ldots, 2^{n}\right\}$ and $\{0,1\}^{n}$ through Collatz, and it's $2^{n}$ periodic.

## 5 Rigorous Proof of Statements

## Theorem 1.

$$
\forall N>0, \forall t \in(0,1), \exists n, 2^{N} \geq n \geq 1,|\mathcal{C}(n)-t| \leq 2^{-N}
$$

Proof. There is a $2^{n}$ periodic bijection for all $n$-digits up to $N$, which is sufficient to claim that modulo $2^{-N}$ the two quantities will be equal.

Proof by Python. Under the assumption that $\arctan (1)=\pi / 4$ and $\arctan (0)=0$
def $f(x)$ :
return $1 /(1+\mathrm{x} * * 2)$
def collatz(x, prec=256, show=False):
$\mathrm{y}=0$
while prec:
if $\mathrm{x} \% 2$ :
$\mathrm{y}+=2$ ** (prec - 257)
$\mathrm{x}=((\mathrm{x} \% 2) *(3 * \mathrm{x}+1)+(\mathrm{x} \% 2==0) * \mathrm{x}) / 2$ prec = prec - 1
if show: print ( $\mathrm{x}, \mathrm{y}$ )
return y
def approx_pi(N, eps=0.0001):
pi = [0]
for $k$ in range (1, N): print(f" pi is approximatively \{4 * sum(pi)\}", end=" \r") pi += [eps * f(collatz(k)) * ((1-eps) ** k)]
print(f" pi is approximatively $\{4 * \operatorname{sum}(p i)\} \backslash t i t e r a t i o n s\{N\} ")$
return sum(pi) * 4
approx_pi(100000)

